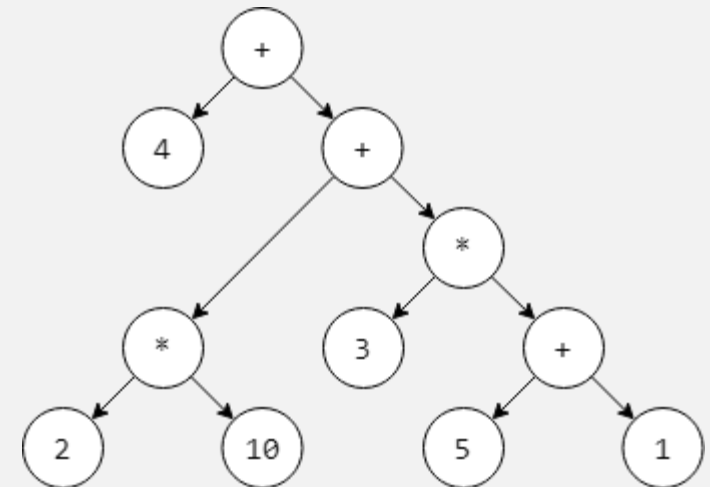


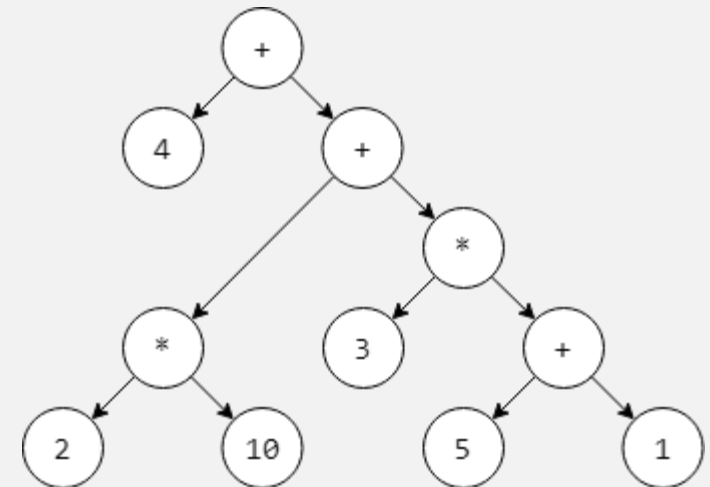
UMass Boston Computer Science  
**CS450 High Level Languages** (section 2)  
**ASTs and Interpreters**

Monday, November 4, 2024



## Logistics

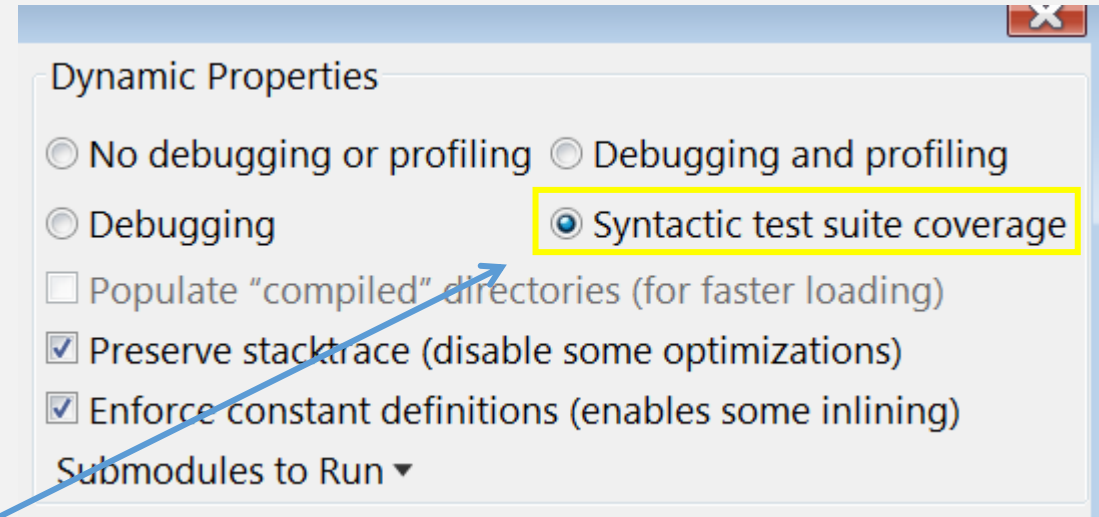
- HW 9 “out”
  - See: in-class work 11/4 and 11/6
  - due: Mon 11/11 12pm (noon) EST
- HW 10 – extension of “hw9”
  - Out: Tue 11/5
  - due: Mon 11/18 12pm (noon) EST
- no lecture: Veteran’s Day Mon 11/11



Previously

# HW Minimum Submission Requirements

- “main” runs without errors
- Tests run without errors
- 100% (Test / Example) “Coverage”
  - In “Choose Language” Menu
  - NOTE: only works with single files



```
;; YCoord is either
;; - before target
;; - in target
;; - after target
;; - out of scene
(define (PENDING-Note? n) (PENDING? (Note-state n)))
(define (HIT-Note? n) (HIT? (Note-state n)))
(define (MISSED-Note? n) (MISSED? (Note-state n)))
(define (OUTOFSCENE-Note? n) (OUTOFSCENE? (Note-state n)))
(define out-Note? OUTOFSCENE-Note?)

;; NEW
;; A WorldState is a List<Note>

(define (num-Notes w) (length w))
```

This code was not run

# HW Minimum Submission Requirements

- “main” runs without errors
- Tests run without errors

Code should never get into a state where this is true!

**Incremental programming!**

# Function Design Recipe

1. **Name**
2. **Signature** – types of the function input(s) and output
3. **Description** – explain (in English prose) the function behavior
4. **Examples** – show (using `rackunit`) the function behavior
5. **Template** – sketch out the function structure (using input's Data Definition)
6. **Code** – implement the rest of the function (arithmetic)
7. **Tests** – check (using `rackunit`) the function behavior

# Incremental Programming

1. **Name**
2. **Signature** – types of the function input(s) and output
3. **Description** – explain (in English prose) the function behavior
4. **Examples** – show (using `rackunit`) the function behavior
5. **Template** – sketch out the function structure (using input's Data Definition)
6. **Code** – implement the rest of the function (arithmetic)  
Start: by filling in with “placeholders”
7. **Tests** – check (using `rackunit`) the function behavior

Code should never  
be crashing!

# Incremental Programming

1. **Name**
2. **Signature** – types of the function input(s) and output
3. **Description** – explain (in English prose) the function behavior
4. **Examples** – show (using `rackunit`) the function behavior
5. **Template** – sketch out the function structure (using input's Data Definition)
6. **Code** – implement the rest of the function (arithmetic)
7. **Tests** – check (using `rackunit`) the function behavior

Code should never be crashing!

This way: always know where the “bug” is

Then: make small code changes and test immediately

Tests (and example tests) should always be passing!

*Previously*

# Intertwined Data Definitions

- Come up with a Data Definition for ...
- ... valid Racket Programs



# Basic Valid Racket Programs

- 1
- “one”
- (+ 1 2)

```
;; A RacketProg is a:
```

```
;; - Number
```

```
;; - String
```

```
;; - ???
```

# Valid Racket Programs

- 1
- “one”
- (+ 1 2)

```
;; A RacketProg is a:  
;; - Atom
```

```
;; - ???
```

```
;; An Atom is one of:  
;; - Number  
;; - String
```

# Valid Racket Programs

• (+ 1 2) ← List of ... atoms?

“symbol”

```
;; A RacketProg is a:
```

```
;; - Atom
```

```
;; - List<Atom> ???
```

```
;; An Atom is one of:
```

```
;; - Number
```

```
;; - String
```

```
;; - Symbol
```

# Valid Racket Programs

- `(* (+ 1 2) (- 4 3))`

Tree?

- `(* (+ 1 2) (- 4 3) (/ 10 5))`

Each tree "node" is a list, of ... RacketProgs ??

But: how many values does each node have?? Unknown!

```
;; A RacketProg is a:
```

```
;; - Atom
```

```
;; - List<Atom> ???
```

```
;; - Tree<???
```

```
;; An Atom is one of:
```

```
;; - Number
```

```
;; - String
```

```
;; - Symbol
```

# Valid Racket Programs

- `(* (+ 1 2) (- 4 3))`

Tree?

- `(* (+ 1 2) (- 4 3) (/ 10 5))`

Each tree "node" is a list, of ... RacketProgs ??

But: how many values does each node have??

```
;; A RacketProg is a:  
;; - Atom  
;; - ProgTree
```

```
;; An Atom is one of:  
;; - Number  
;; - String  
;; - Symbol
```

```
;; A ProgTree is one of:  
;; - empty  
;; - (cons RacketProg ProgTree)
```

Recursive Data Def!

# Valid Racket Programs

Also, **Intertwined Data Defs!**

```
;; A RacketProg is a:  
;; - Atom  
;; - ProgTree
```

```
;; A ProgTree is one of:  
;; - empty  
;; - (cons RacketProg ProgTree)
```

```
;; An Atom is one of:  
;; - Number  
;; - String  
;; - Symbol
```

# Intertwined Data

- A set of Data Definitions that reference each other
- Templates should be defined together ...

```
;; A RacketProg is a:  
;; - Atom  
;; - ProgTree
```

```
;; A ProgTree is one of:  
;; - empty  
;; - (cons RacketProg ProgTree)
```

```
;; An Atom is one of:  
;; - Number  
;; - String  
;; - Symbol
```

# Intertwined Data

- A set of Data Definitions that reference each other
- Templates should be defined together ...
  - ... and should reference each other's templates (when needed)

```
;; A RacketProg is one of:  
;; - Atom  
;; - ProgTree
```

```
(define (prog-fn p) ...)
```

```
;; A ProgTree is one of:  
;; - empty  
;; - (cons RacketProg ProgTree)
```

```
(define (ptree-fn t) ...)
```

```
;; An Atom is one of:  
;; - Number  
;; - String  
;; - Symbol
```

```
(define (atom-fn a) ...)
```

???



# Intertwined Templates

```
;; A RacketProg is one of:  
;; - Atom  
;; - ProgTree
```

```
(define (prog-fn s)  
  (cond  
    [(atom? s) ... (atom-fn s) ...]  
    [else ... (ptree-fn s) ...]))
```

```
;; A ProgTree is one of:  
;; - empty  
;; - (cons RacketProg ProgTree)
```

```
(define (ptree-fn t)  
  (cond  
    [(empty? t) ...]  
    [else ... (prog-fn (first t)) ... (ptree-fn (rest t)) ...]))
```

```
;; An Atom is one of:  
;; - Number  
;; - String  
;; - Symbol
```

```
(define (atom-fn a)  
  (cond  
    [(number? a) ... ]  
    [(string? a) ... ]  
    [else ... ]))
```

**Intertwined data have  
intertwined templates!**

# A “Racket Prog” = S-expression!

```
;; A RacketProg Sexpr is one of:  
;; - Atom  
;; - ProgTree
```

```
(define (sexpr-fn s)  
  (cond  
    [(atom? s) ... (atom-fn s) ...]  
    [else ... (ptree-fn s) ...]))
```

```
;; A ProgTree is one of:  
;; - empty  
;; - (cons RacketProg Sexpr ProgTree)
```

```
(define (ptree-fn t)  
  (cond  
    [(empty? t) ...]  
    [else ... (sexpr-fn (first t)) ... (ptree-fn (rest t)) ...]))
```

```
;; An Atom is one of:  
;; - Number  
;; - String  
;; - Symbol
```

```
(define (atom-fn a)  
  (cond  
    [(number? a) ... ]  
    [(string? a) ... ]  
    [else ... ]))
```

# Counting Symbols

```
;; count : Symbol Sexpr -> Nat  
;; Computes the number of times the given  
;; symbol appears in the given s-expression
```

```
(define (count sym se)  
  (cond  
    [(atom? s) ... (atom-fn s) ...]  
    [else ... (ptree-fn s) ...]))
```

```
;; count-ptree : Symbol ProgTree -> Nat
```

```
(define (count-ptree sym pt)  
  (cond  
    [(empty? pt) ...]  
    [else ... (sexpr-fn (first pt)) ... (ptree-fn (rest pt)) ...]))
```

```
;; count-atom : Symbol Atom -> Nat
```

```
(define (count-atom sym a)  
  (cond  
    [(number? a) ... ]  
    [(string? a) ... ]  
    [else ... ]))
```

# Counting Symbols

```
;; count : Symbol Sexpr -> Nat  
;; Computes the number of times the given  
;; symbol appears in the given s-expression
```

```
(define (count sym se)  
  (cond  
    [(atom? s) (count-atom sym se)]  
    [else (count-ptree sym se)]))
```

```
;; count-ptree : Symbol ProgTree -> Nat
```

```
(define (count-ptree sym pt)  
  (cond  
    [(empty? pt) ...]  
    [else ... (sexpr-fn (first pt)) ... (ptree-fn (rest pt)) ...]))
```

```
;; count-atom : Symbol Atom -> Nat
```

```
(define (count-atom sym a)  
  (cond  
    [(number? a) ... ]  
    [(string? a) ... ]  
    [else ... ]))
```

# Counting Symbols

```
;; count : Symbol Sexpr -> Nat  
;; Computes the number of times the given  
;; symbol appears in the given s-expression
```

```
(define (count sym se)  
  (cond  
    [(atom? s) (count-atom sym se)]  
    [else (count-ptree sym se)]))
```

```
;; count-ptree : Symbol ProgTree -> Nat
```

```
(define (count-ptree sym pt)  
  (cond  
    [(empty? pt) ...]  
    [else ... (sexpr-fn (first pt)) ... (ptree-fn (rest pt)) ...]))
```

```
;; count-atom : Symbol Atom -> Nat
```

```
(define (count-atom sym a)  
  (cond  
    [(symbol? a)  
     (if (symbol=? sym a) 1 0)]  
    [else 0]))
```

# Counting Symbols

```
;; count : Symbol Sexpr -> Nat  
;; Computes the number of times the given  
;; symbol appears in the given s-expression
```

```
(define (count sym se)  
  (cond  
    [(atom? s) (count-atom sym se)]  
    [else (count-ptree sym se)]))
```

```
;; count-ptree : Symbol ProgTree -> Nat
```

```
(define (count-ptree sym pt)  
  (cond  
    [(empty? pt) 0]  
    [else ... (sexpr-fn (first pt)) ... (ptree-fn (rest pt)) ...]))
```

```
;; count-atom : Symbol Atom -> Nat
```

```
(define (count-atom sym a)  
  (cond  
    [(symbol? a)  
     (if (symbol=? sym a) 1 0)]  
    [else 0]))
```

# Counting Symbols

```
;; count : Symbol Sexpr -> Nat  
;; Computes the number of times the given  
;; symbol appears in the given s-expression
```

```
(define (count sym se)  
  (cond  
    [(atom? s) (count-atom sym se)]  
    [else (count-ptree sym se)]))
```

```
;; count-ptree : Symbol ProgTree -> Nat
```

```
(define (count-ptree sym pt)  
  (cond  
    [(empty? pt) 0]  
    [else (+ (count sym (first pt))  
             (count-ptree sym (rest pt)))]))
```

```
;; count-atom : Symbol Atom -> Nat
```

```
(define (count-atom sym a)  
  (cond  
    [(symbol? a)  
     (if (symbol=? sym a) 1 0)]  
    [else 0]))
```

# Syntax vs Semantics (Spoken Language)

## Syntax

- Specifies: validity of language structures
  - E.g., sentence = noun (subject) + verb + noun (object)
- “the ball threw the child”
  - Syntactically: valid!
  - Semantically: ???

## Semantics

- Specifies: meaning of language structures



# Syntax vs Semantics (Programming Language)

## **Syntax**

- Specifies: validity of language structures
  - E.g., ???

## **Semantics**

- Specifies: meaning of language structures

# Syntax vs Semantics (Programming Language)

## Syntax

- Specifies: validity of ~~language structures~~ **programs!**
  - E.g., valid Racket program = s-expressions
  - E.g., valid Python program = ...

## Semantics

- Specifies: meaning of ~~language structures~~ **programs!**

**Q:** What is the “meaning” of a program?

**A:** The result from “running” it

How does a program “run”?

# Running Programs: `eval`

```
;; eval : Sexpr -> Result  
;; “runs” a given Racket program, producing a “result”
```

An “eval” function turns a “program” into a “result”

An “eval” function is more generally called an **interpreter**

(Programs are usually not directly interpreted)

More commonly, a high-level program is first **compiled** to a lower-level language (and then interpreted)

**Q:** What is the “meaning” of a program?

**A:** The result from “running” it

How does a program “run”?

From  
Lecture 1

“high” level  
(easier for humans  
to understand)

NOTE: This hierarchy is approximate

“declarative”

More commonly, a  
high-level program  
is first **compiled** to  
a lower-level  
language (and then  
intrepreted)

“imperative”

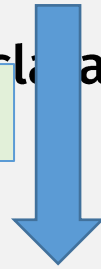
“low” level  
(runs on cpu)

English	
Specification langs	Types? pre/post cond?
Markup (html, markdown)	tags
Database (SQL)	queries
Logic Program (Prolog)	relations
Lazy lang (Haskell, R)	Delayed computation
Functional lang (Racket)	Expressions (no stmts)
JavaScript, Python	“eval”
C# / Java	GC (no alloc, ptrs)
C++	Classes, objects
C	Scoped vars, fns
Assembly Language	Named instructions
Machine code	Binary

“high” level  
(easier for humans  
to understand)

**surface language**

“declarative”  
**compiler**



**target language**

“imperative”

“low” level  
(runs on cpu)

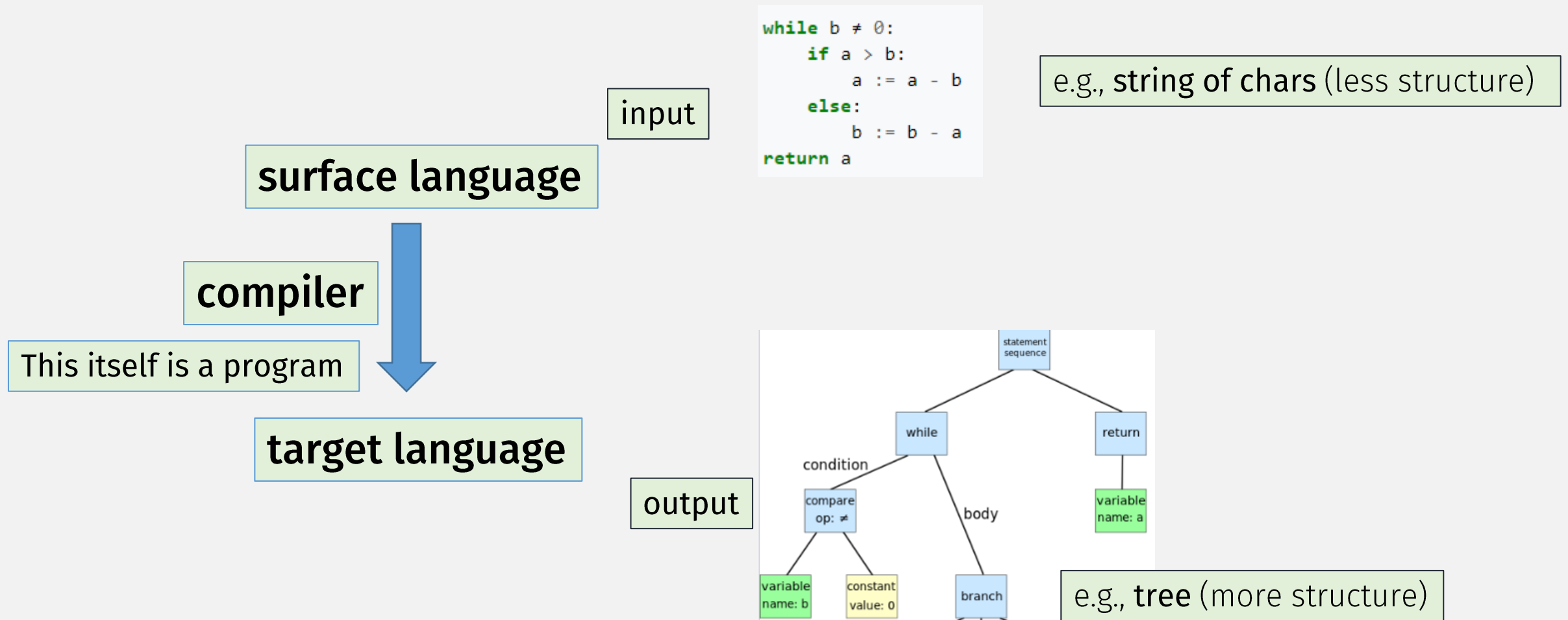
Specification langs
Markup (html, markdown)
Database (SQL)
Logic Program (Prolog)
Lazy lang (Haskell, R)
Functional lang (Racket)
JavaScript, Python
C# / Java
C++
C
Assembly Language
Machine code

Common **target** languages:

- bytecode (e.g., JS, Java)
- assembly
- machine code

A **virtual machine** is just a  
**bytecode interpreter**

A (hardware) **CPU** is just a  
**machine code interpreter!**



## Semantics

- Specifies: meaning of language structures
- So: to “run” a program, we need to see the structure first

```
while b ≠ 0:
  if a > b:
    a := a - b
  else:
    b := b - a
return a
```

e.g., string of chars (less structure)

input

surface language

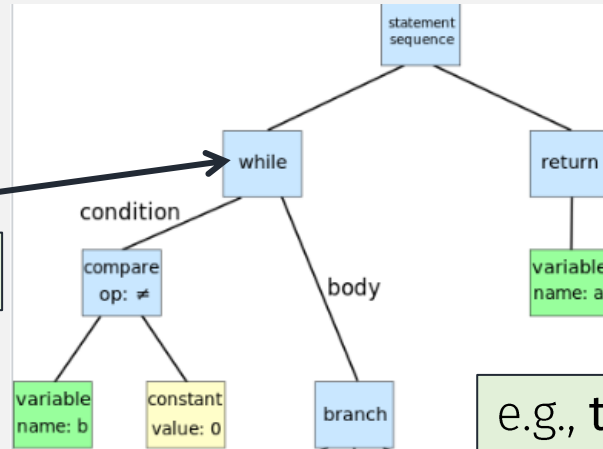
Compiler, step 1

= parser

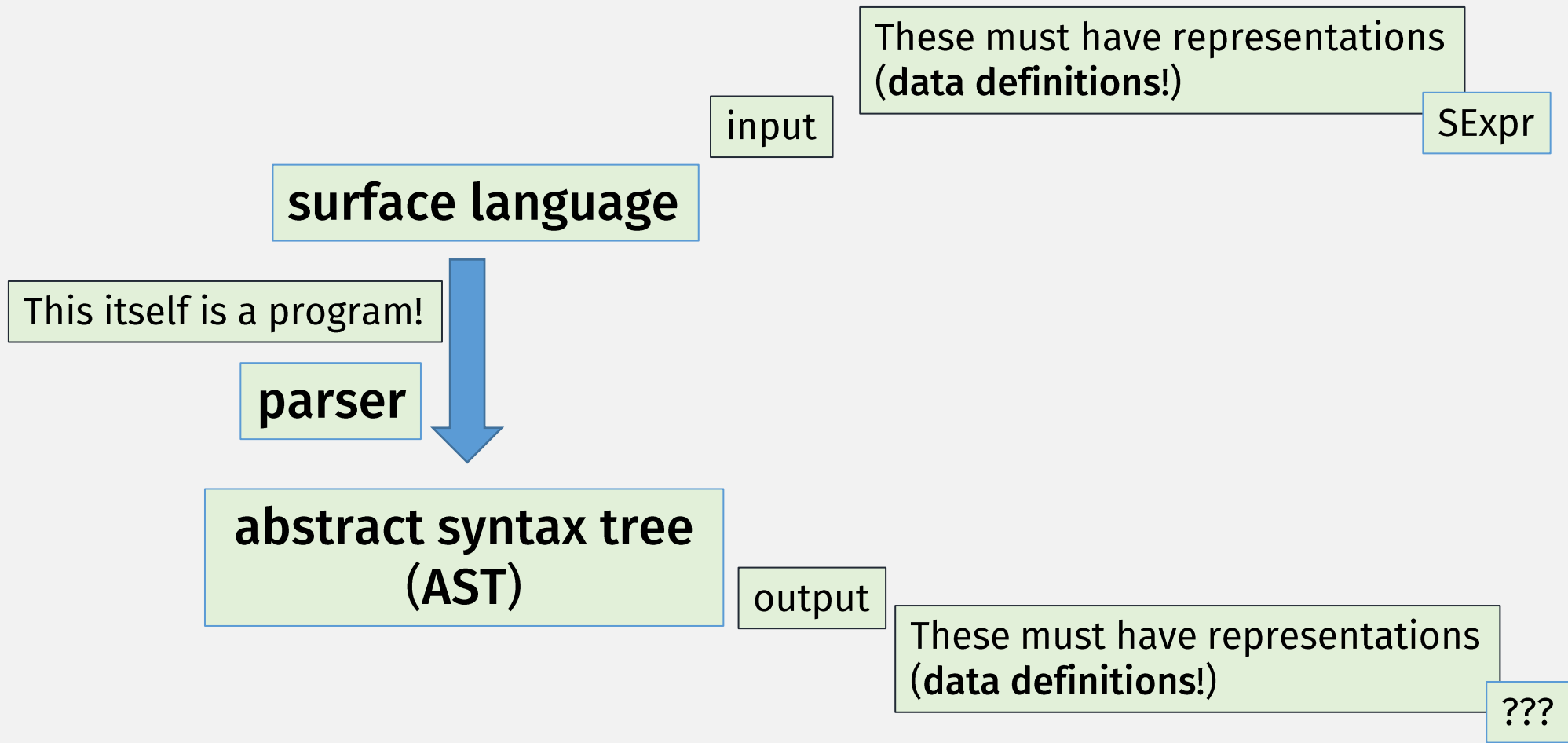
(a compiler actually has many steps... take a compilers course!)

abstract syntax tree (AST)

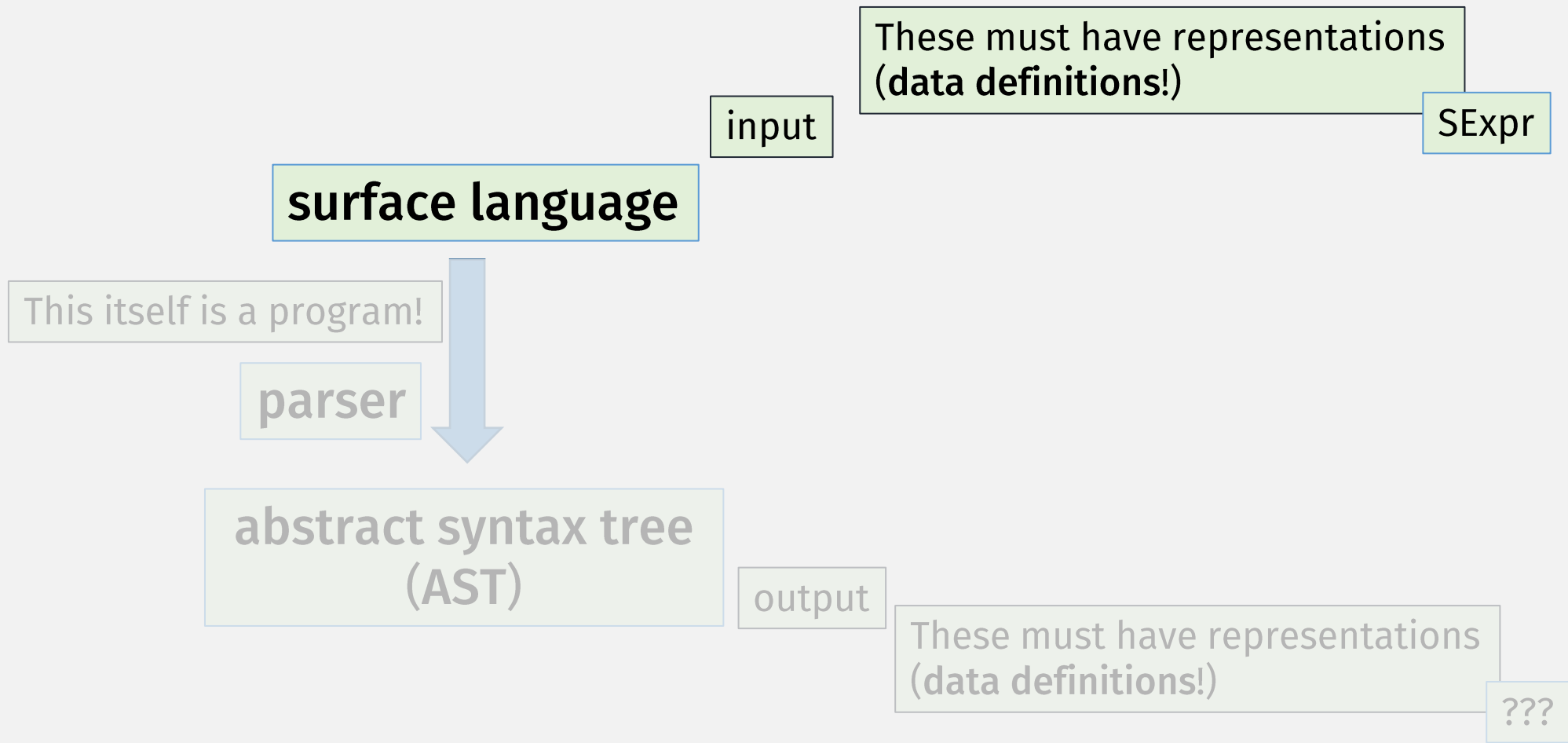
output



e.g., tree (more structure)







surface language

input

These must have representations  
(data definitions!)

SExpr

```
;; A SimpleSexpr (Ssexpr) is one of:  
;; - Number  
;; - (list '+ Ssexpr Ssexpr)  
;; - (list '- Ssexpr Ssexpr)
```

# Data Definition Template

When a **Data Definition** is an **itemization** of compound data ...

- **Template** =
  - cond to distinguish cases
  - “Getters” to extract pieces
  - recursive calls

```
;; A SimpleSexpr (Ssexpr) is one of:  
;; - Number  
;; - (list '+ Ssexpr Ssexpr)  
;; - (list '- Ssexpr Ssexpr)
```

```
(define (ss-fn s)  
  (cond  
    [(number? s) ... ]  
    [(and (list? s) (equal? '+ (first s)))  
     ... (ss-fn (second s)) ... (ss-fn (third s)) ... ]  
    [(and (list? s) (equal? '- (first s)))  
     ... (ss-fn (second s)) ... (ss-fn (third s)) ... ]))
```

Cond guards must distinguish the different cases

“getters”

Recursive call(s)

# Interlude: pattern matching (again)

When a **Data Definition** is an **itemization** of compound data ...

- **Template** =
  - ~~cond to distinguish cases~~
  - **match** = cond + getters
  - recursive calls

```
;; A SimpleSexpr (Ssexpr) is one of:  
;; - Number  
;; - (list '+ Ssexpr Ssexpr)  
;; - (list '- Ssexpr Ssexpr)
```

```
(define (ss-fn s)  
  (match s  
    [(? number?) ... ]  
    [(+ ,x ,y) "Quasiquote" pattern  
     ... (ss-fn x) ... (ss-fn y) ... ]  
    [ `(- ,x ,y) Symbols match exactly  
     ... (ss-fn x) ... (ss-fn y) ... ]))
```

Match patterns

Predicate pattern

"Quasiquote" pattern

Symbols match exactly

"Unquote" defines new variable name (for value at that position)

# Interlude: pattern matching (again)

- See Racket docs for the full pattern language

The grammar of *pat* is as follows, where non-italicized identifiers are recognized symbolically (i.e., not by binding).

<code>pat ::= id</code>	match anything, bind identifier
<code>  (id datum)</code>	match anything, bind identifier
<code>  _</code>	match anything
<code>  literal</code>	match literal
<code>  (quote datum)</code>	match <code>equal?</code> value
<code>  (list lvp ...)</code>	match sequence of <i>lvps</i>
<code>  (list-rest lvp ... pat)</code>	match <i>lvps</i> consed onto a <i>pat</i>
<code>  (list* lvp ... pat)</code>	match <i>lvps</i> consed onto a <i>pat</i>
<code>  (list-no-order pat ...)</code>	match <i>pats</i> in any order
<code>  (list-no-order pat ... lvp)</code>	match <i>pats</i> in any order
<code>  (vector lvp ...)</code>	match vector of <i>pats</i>
<code>  (hash-table (pat pat) ...)</code>	match hash table
<code>  (hash-table (pat pat) ...+ ooo)</code>	match hash table
<code>  (cons pat pat)</code>	match pair of <i>pats</i>
<code>  (mcons pat pat)</code>	match mutable pair of <i>pats</i>
<code>  (box pat)</code>	match boxed <i>pat</i>
<code>  (struct-id pat ...)</code>	match <i>struct-id</i> instance
<code>  (struct struct-id (pat ...))</code>	match <i>struct-id</i> instance
<code>  (regexp rx-expr)</code>	match string
<code>  (regexp rx-expr pat)</code>	match string, result with <i>pat</i>
<code>  (pregexp px-expr)</code>	match string
<code>  (pregexp px-expr pat)</code>	match string, result with <i>pat</i>
<code>  (and pat ...)</code>	match when all <i>pats</i> match
<code>  (or pat ...)</code>	match when any <i>pat</i> match
<code>  (not pat ...)</code>	match when no <i>pat</i> matches
<code>  (app expr pats ...)</code>	match ( <i>expr</i> value) output values to <i>pats</i>
<code>  (? expr pat ...)</code>	match if ( <i>expr</i> value) and <i>pats</i>
<code>  (quasiquote qp)</code>	match a quasipattern
<code>  derived-pattern</code>	match using extension

# Interlude: pattern matching (again)

When a **Data Definition** is an itemization of compound data ...

- **Template =**
  - ~~cond~~ to distinguish cases
  - match = cond + getters
  - recursive calls

match can be more concise and readable

```
(define (ss-fn s)
  (match s
    [(? number?) ... ]
    [`(+ ,x ,y)
     ... (ss-fn x) ... (ss-fn y) ... ]
    [`(- ,x ,y)
     ... (ss-fn x) ... (ss-fn y) ... ])))
```

VS

```
(define (ss-fn s)
  (cond
    [(number? s) ... ]
    [(and (list? s) (equal? '+ (first s)))
     ... (ss-fn (second s)) ...
     ... (ss-fn (third s)) ... ]
    [(and (list? s) (equal? '- (first s)))
     ... (ss-fn (second s)) ...
     ... (ss-fn (third s)) ... ])))
```

surface language

input

These must have representations  
(data definitions!)

SExpr

This itself is a program!

parser

abstract syntax tree  
(AST)

output

```
;; A SimpleSexpr (Ssexpr) is one of:  
;; - Number  
;; - (list '+ Ssexpr Ssexpr)  
;; - (list '- Ssexpr Ssexpr)
```

These must have representations  
(data definitions!)

???

surface language

This itself is a program!

parser

abstract syntax tree  
(AST)

These must have representations

```
;; An AST is one of:  
;; - (num Number)  
;; - (plus AST AST)  
;; - (minus AST AST)  
;; Interp: Tree structure for Ssexpr prog  
(struct num [val])  
(struct plus [left right])  
(struct minus [left right])
```

output

These must have representations  
(data definitions!)

???



```
;; An AST is one of:  
;; - (num Number)  
;; - (plus AST AST)  
;; - (minus AST AST)  
;; Interp: Tree structure for Ssexpr prog  
(struct num [val])  
(struct plus [left right])  
(struct minus [left right])
```

• **Template =**

```
(define (ast-fn p)  
  (cond  
    [(num? p) ... ]  
    [(plus? p) ... (ast-fn (plus-left p))  
                  ... (ast-fn (plus-right p)) ... ]  
    [(minus? p) ... (ast-fn (minus-left p))  
                  ... (ast-fn (minus-right p)) ... ]))
```

• **Template** (with match) =

```
;; An AST is one of:  
;; - (num Number)  
;; - (plus AST AST)  
;; - (minus AST AST)  
;; Interp: Tree structure for Ssexpr prog  
(struct num [val])  
(struct plus [left right])  
(struct minus [left right])
```

```
(define (ast-fn p)  
  (cond match p  
    [(num n) ... ]  
    [(plus x y) ... (ast-fn x) ...  
     ... (ast-fn y) ... ]  
    [(minus x y) ... (ast-fn x) ...  
     (ast-fn y) ... ]))
```

Struct name

Struct patterns

Extracts and names fields

# In-class Coding 11/4 #1 (HW9): parser

```
;; parse: SimpleSexpr -> AST
;; Converts a (simple) S-expression to language AST
```

```
;; A SimpleSexpr (Ssexpr) is a:
;; - Number
;; - (list '+ Ssexpr Ssexpr)
;; - (list '- Ssexpr Ssexpr)
```

```
;; An AST is one of:
;; - (num Number)
;; - (plus AST AST)
;; - (minus AST AST)
;; Interp: Tree structure for Ssexpr
(struct num [val])
(struct plus [left right])
(struct minus [left right])
```

# In-class Coding 11/4 #2 (HW9): run

```
;; run: AST -> Result  
;; computes the result of given program AST
```

```
;; A Result is a ... ???
```

```
;; An AST is one of:  
;; - (num Number)  
;; - (plus AST AST)  
;; - (minus AST AST)  
;; Interp: Tree structure for Ssexpr  
(struct num [val])  
(struct plus [left right])  
(struct minus [left right])
```

```
;; eval-ssexpr : Ssexpr -> Result  
(define eval-ssexpr  
  (compose run parse))
```